

Peer-to-Peer Traffic Localization as a Service

Moritz Steiner and Matteo Varvello
Bell Labs, Alcatel-Lucent, USA
{*first.last*}@alcatel-lucent.com

Abstract—Peer-to-Peer (P2P) file-sharing applications generate extremely high network transit costs for the ISPs due to their network agnostic nature. A number of strategies have been deployed to localize P2P traffic, mostly relying on the communication between the peers and the directory servers (trackers). However, recently many users favor distributed tracking based upon Distributed Hash Tables (DHTs). In this demonstration, we present the first traffic localization mechanism for DHT-based file-sharing networks. It runs as a service in the cloud. Thus, it does not require any equipment installation at an ISP, nor cooperation between an ISP and P2P networks.

I. INTRODUCTION

BitTorrent is by far the most popular Peer-to-Peer (P2P) protocol, adopted by several file-sharing applications such as μ Torrent [10] and Azureus [3]. The BitTorrent protocol ignores the location of peers in the network when determining how files are exchanged, thus generating a large volume of expensive inter-ISP traffic. In the past, tracker-based traffic localization strategies have shown they effectively reduce BitTorrent inter-ISPs traffic [13].

Recently, BitTorrent introduced a distributed tracking feature: peers can now locate which peers hold a copy or a portion of a file by querying a Distributed Hash Table (DHT). This feature, originally thought as a backup solution for the trackers, is now supported by most BitTorrent client implementations and it attracts a large share of BitTorrent user requests [11]. The addition of a DHT to the BitTorrent network decreases from the effectiveness of tracker-based traffic localization mechanisms.

In [11], we designed and prototyped the first traffic localization mechanism for DHT-based file-sharing networks. Our prototype targets the BitTorrent Mainline DHT [6] and can be run as a service in the cloud. No equipment installation is required at an ISP, nor does it require cooperation between an ISP and P2P networks (which is unlikely to occur). This demonstration shows the prototype used in the experimental evaluation of [11].

II. DESIGN OVERVIEW

Our localization mechanism works in two steps. First, we intercept all the messages from peers announcing in the DHT that they hold a file or a portion of it. Then, we intercept all the requests for this file and answer with *local* peer-sets, i.e., sets of peers located at the same ISPs as the requesting peers. We now describe both steps in detail.

A single entity can join a P2P network many times with many distinct logical identities called *sybils* [4]. To intercept announces and requests for a file, we insert several sybils in the

DHT with nodeIDs close to the `info_hash` of the file [9]. The sybils are constantly aware of the peers that hold the file as well as the peers requesting it. Under this premise, localization is straightforward. Our sybils simply respond to the queries for this file with localized peer-sets. If only a few local peers are available, external peers are used to complete the peer-set.

Our localization mechanism targets only popular files for two reasons. First, only the traffic associated with files requested by more than one peer from the same ISP at the same time has potential for localization. Second, we aim to minimize the number of files to be localized. In fact, the number of sybils we need to insert in the DHT to achieve traffic localization scales linearly with the number of files that we localize.

III. MAINLINE IMPLEMENTATION

As a proof of concept, we implemented a prototype of our localization mechanism for the Mainline DHT. We pick the Mainline DHT since it has the largest user base with more than eight million active users at any point in time. Note that our localization mechanism can easily be ported to any other DHT-based P2P network build upon the Kademlia protocol [8], e.g., Azureus [3], Emule [5], and Ares [2].

We now briefly overview the relevant messages used in the Mainline DHT. For more details, the interested reader is referred to [6].

`get_peers(I)` – it performs a lookup operation, i.e., retrieve the list of peers holding a copy or a portion of a file with `info_hash I`. The lookup works iteratively. At each intermediary hop of the iteration, peers respond to a `get_peers` message with the IP addresses, ports, and nodeIDs of the peers closest to I . At the final hop of the iteration, peers return the IP addresses and ports of the peers that hold a copy or a portion of the file with `info_hash I`.

`announce_peer(I)` – it is used by a peer P to announce that it hosts a file with `info_hash I`. P sends several `announce_peer` messages to the peers in the DHT whose identifiers are the closest to I (previously retrieved using a `get_peers` message).

`ping(P)` – it checks that peer P is alive.

We now focus on the localization of the traffic associated with a single file identified by `info_hash I`.

In the first step of the localization mechanism, we intercept the `announce_peer` messages for I . To do so, we insert sybils in the DHT with nodeIDs closer to I than any real peer. The insertion of the sybils in the DHT consists of informing real peers whose nodeIDs are close to I about the existence

